

UNIVERSITY OF TWENTE.

Inullen in blokletters/To be completed by student

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1) a) ρ is fraction of time server is working
amount of work arriving to system per time unit = λ EB

b) see reader p 62:

$$\text{PASTA : } p_n = a_n$$

$d_n = a_n$ as number of transitions $n \rightarrow n+1$ per time unit
must equal number of transitions $n+1 \rightarrow n$ per time unit

$$\text{Hence } d_n = p_n$$

c) see reader p 64 for derivation of pgf $P_L(z)$
and use b) to conclude that $P_L(z) = P_{L,d}(z)$

d) is ex 45

$$\tilde{B}(s) = \frac{4}{4+s} \cdot \frac{3}{3+s}$$

e) note: service process of d, $\lambda = 1$, $\rho = \lambda \text{EB} = \frac{7}{12}$

insert $\tilde{B}(1-z)$ from d) in c):

$$P_L(z) = \frac{5}{(z-6)(z-2)}$$

$$\text{EB} = \left. \frac{d}{dz} P_L(z) \right|_{z=1} = \left. \frac{-5(z-8)}{(z-6)^2(z-2)^2} \right|_{z=1} = \frac{-5 \cdot (-6)}{25} = \frac{6}{5}$$

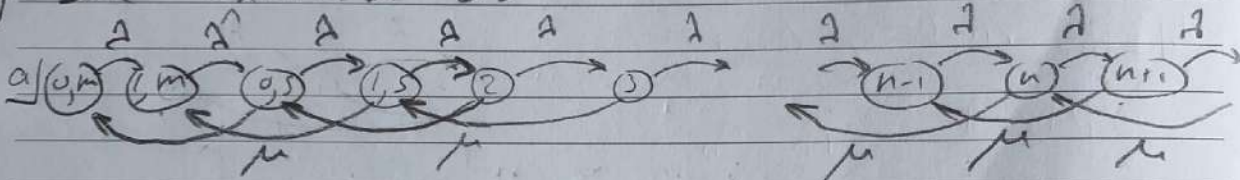
2) see Section 8.3

$$\left. \begin{aligned} ES &= EL^a \frac{1}{\mu} + \frac{1}{\mu} \\ EL^a &= \sum_{n=0}^{\infty} n a_n = \frac{\rho}{1-\rho} \end{aligned} \right\} ES = \frac{1}{(1-\rho)\mu}$$

(P.116)

$$\Rightarrow EL = \frac{\rho}{1-\rho}$$

3) see exercise 55



$n = \#$ waiting for Perry, (n, m) n waiting, Perry moved
 $\lambda = 15 / \text{hour}$, $(n, 0)$ n waiting, Perry moving
 $\mu = 20 / \text{hour}$, $(n, 1)$ n waiting, Perry idle
 $n = 0, 1$

b) re-label states q_k , $k=0, 1, 2, \dots$
 balance equations

$$\left. \begin{aligned} q_0 \lambda &= q_1 \mu \\ q_1 \lambda &= q_0 \lambda + q_2 \mu \\ q_n (\lambda + \mu) &= q_{n-1} \lambda + q_{n+1} \mu \quad n \geq 2 \\ \text{steady state solution} \quad q_n &= \lambda^n \quad n \geq 2 \\ \Rightarrow (\lambda - 1) \left(\lambda - \frac{1}{2} \right) \left(\lambda + \frac{3}{2} \right) &= 0 \Rightarrow \lambda = \frac{1}{2} \\ q_n &= c \lambda^n \quad n \geq 2 \\ \text{boundary eq: } q_0 &= c \frac{\mu}{\lambda} \lambda^2 \\ q_1 &= c \frac{\mu}{\lambda} \lambda^1, c \frac{\mu}{\lambda} \lambda^2 \end{aligned} \right\}$$

normalization $\sum_{n=0}^{\infty} q_n = 1 \Rightarrow C = \frac{3}{4}$

so $P_n = \text{prob } n \text{ cars waiting:}$

$$P_n = q_{nr2} = \frac{3}{4} \left(\frac{1}{2}\right)^{nr2} \quad n \geq 2$$

$$P_1 = q_1 + q_3 = \frac{3}{8} + \frac{3}{4} \left(\frac{1}{2}\right)^3$$

$$P_2 = q_0 + q_2 = \frac{1}{4} + \frac{3}{4} \left(\frac{1}{2}\right)^2$$

4/ a) This is example 4.2.3 page 55 of reader

b) idem

c) This is proof of Theorem 4.4.1
for $f(l, n) = 1(l=n)$, $S(l, n) = 1(l \leq n)$

and using the rules of the time-reversed queue from ex 4.2.3

5/ correction: jobs served at queue 2 route to queue 3

a) $\bar{n} = (n_1, n_2, n_3)$ with $n_i = \#$ customers in queue i

$$\mathcal{S} = \{ \bar{n} : n_i \geq 0, i=1,2,3 \}$$

$$q(n, n+e_1) = \mu_0$$

$$q(n, n-e_1+e_2) = \frac{1}{3}\mu_1$$

$$q(n, n-e_1+e_3) = \frac{2}{3}\mu_1$$

$$q(n, n-e_2+e_1) = \mu_2$$

$$q(n, n-e_3) = \mu_3$$

b) traffic equations $\lambda_j = \mu_0 p_{0j} + \sum_{i=1}^3 \lambda_i p_{ij} \quad j=1,2,3$

$$\begin{aligned} \lambda_1 &= \mu_0 \\ \lambda_2 &= \frac{1}{3} \lambda_1 \\ \lambda_3 &= \frac{2}{3} \lambda_1 + \lambda_2 \end{aligned}$$

solution: $\begin{aligned} \lambda_1 &= \mu_0 \\ \lambda_2 &= \frac{1}{3} \mu_0 \\ \lambda_3 &= \frac{2}{3} \mu_0 \end{aligned}$ (obvious as all customers arrive at queue 1)

stability condition $\lambda_i < \mu_i \quad i=1,2,3$

so $\mu_0 < \mu_1, \mu_0 < 3\mu_2, \mu_0 < \mu_3$

stab condition: $\mu_0 < \min\{\mu_1, 3\mu_2, \mu_3\}$

c) $\pi(\vec{n}) = \prod_{i=1}^3 (1-\rho_i) \rho_i^{n_i} \quad n_i \geq 0$ with $\rho_i = \lambda_i / \mu_i$

The network is a feed forward network
so we can indeed prove this result using Burke's theorem

d) How $n_3 \leq M_3$ customers arriving to queue 3 will jump over queue 3 if $n_3 = M_3$

$$\pi(\vec{n}) = (1-\rho_1) \rho_1^{n_1} (1-\rho_2) \rho_2^{n_2} \frac{1-\rho_3}{1-\rho_3^{M_3+1}} \rho_3^{n_3} \quad \begin{aligned} n_i &\geq 0 \\ n_2 &\geq 0 \\ 0 \leq n_3 \leq M_3 \end{aligned}$$

proof: either use partial balance or Burke's theorem as the network remains feed forward and queue 3 is now a $M/M/1/M_3$ queue.