

Exam Limits to Computing (201300042)

Thursday, November 3, 2022, 8:45 – 11:45

- You can bring five sheets of A4 paper with arbitrary content to the exam.
- Electronic devices of any kind, books, ... are not allowed.
- This exam consists of four problems.
- Please start a new page for each problem.
- The total number of points is 60. Sufficient for passing are 30 points.

1. Computability

Let

$$FP = \{g \in \mathcal{G} \mid \varphi_g(g) = g\}.$$

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- (a) (8 points) Is $FP \in RE$? Prove your answer.
- u* (b) (8 points) Is $FP \in REC$? Prove your answer.
- (c) (2 points) Is $FP \in co-RE$? Prove your answer.
- (d) (5 points) Prove that FP is not an index set.

2. NP-Completeness

An instance of SETCOVER is a finite set U together with subsets $S_1, \dots, S_m \subseteq U$ and a $k \in \mathbb{N}$. The question is if there exists k of these subsets that cover all elements in U .

More formally, an instance as described above is a "yes" instance if there exists a set $I \subseteq \{1, \dots, m\}$ with $|I| \leq k$ such that

$$\bigcup_{i \in I} S_i = U.$$

(13 points) Prove that SETCOVER is NP-complete.

Hint: VERTEXCOVER = $\{(G, k) \mid \text{undirected graph } G \text{ has a vertex cover of size } k\}$ is NP-complete.

3. Complexity

Let $E = \text{DTime}(2^{O(n)})$. Recall that $\text{EXP} = \bigcup_{c>0} \text{DTime}(2^{O(n^c)})$.

- (a) (2 points) Prove that $E \subsetneq \text{EXP}$.
- (b) (4 points) Prove that E is not closed under polynomial-time many-one reductions. This means that there are problems $A \notin E$ and $B \in E$ with $A \leq_P B$.

Hint: Consider some $A \in \text{EXP} \setminus E$ and $B = \{x\#1^{f(|x|)} \mid x \in A\}$ for some well-chosen function f .

- (c) (2 points) Prove that $E \neq \text{PSPACE}$.

Hint: You cannot conclude $E \subsetneq \text{PSPACE}$ or $\text{PSPACE} \subsetneq E$ from this. You can only conclude that the two classes are not equal.

4. Questions

Are the following statements true or false? Justify your answers briefly.

- (a) (2 points) $\text{co-NSpace}(n^2) \subsetneq \text{DSpace}(n^5)$.
- (b) (2 points) For all $L \subseteq \mathbb{N}$, the following holds: If $L \in \text{REC}$, then χ_L is total.
- (c) (2 points) For all $L \subseteq \mathbb{N}$, the following holds: If χ_L is total, then $L \in \text{REC}$.
- (d) (2 points) If $L \leq H_0$, then L is recursively enumerable.
- (e) (2 points) If $P = \text{NP}$, then 2SAT is NP-complete.
- (f) (2 points) If 2SAT is NP-complete, then $P = \text{NP}$.

For the last two parts of this exercise, let $\text{QP} = \text{NP} \cup \text{co-NP}$.

- (g) (2 points) QP is closed under polynomial-time many-one reductions.
- (h) (2 points) If there exists a QP-complete problem, then $\text{NP} = \text{co-NP}$.