

Resit Linear Optimization

Friday, February 1, 2019, 13.45-16.45

Give motivations for all your answers. You are not allowed to use a calculator.

1. We can produce two different products a and b , using two resources, r and s .
 One unit of product a requires r_a units of resource r and s_a units of resource s and yields p_a dollars.
 One unit of product b requires r_b units of resource r and s_b units of resource s and yields p_b dollars.
 We buy our resources from a supplier, who charges us c_r dollars per unit of resource r and c_s dollars per unit of resource s .
 Due to demand, we can sell at most d_a units of product a and d_b units of product b .

- a) Construct a linear program that maximizes our profit (yield minus cost).
- b) Suppose that our supplier offers us 10% discount if we spend at least t dollars on resources. Note that we can choose whether or not to make use of this offer. Explain how you would find the maximum profit in this case.
 Hint: Construct a new linear program, and use it together with the linear program from part a).

2. a) Prove that the set $\{(x, y) \in \mathbb{R}^2 \text{ for which } -|x| \leq y\}$ is not convex.

- b) Prove that the set $\{(x, y) \in \mathbb{R}^2 \text{ for which } y \leq -|x|\}$ is convex.

3. a) Prove or give a counterexample:
 For all polyhedra $P \subseteq \mathbb{R}^2$, for all vertices $\mathbf{x}^* \in P$, there exists a line $l = \{\mathbf{x} \in \mathbb{R}^2 \mid \mathbf{a}'\mathbf{x} = \mathbf{b}\}$ for which $l \cap P = \mathbf{x}^*$, i.e. \mathbf{x}^* is the only element of P on l .
- b) Prove or give a counterexample:
 For all polyhedra $P \subseteq \mathbb{R}^2$, for all feasible solutions $\mathbf{x}^* \in P$, the following holds:
 If there exists a line $l = \{\mathbf{x} \in \mathbb{R}^2 \mid \mathbf{a}'\mathbf{x} = \mathbf{b}\}$ for which $l \cap P = \mathbf{x}^*$, then \mathbf{x}^* is a vertex of P .

4. Solve the following linear program using the simplex method.

$$\begin{aligned} \min \quad & -4x_1 + 2x_2 \\ \text{s.t.} \quad & 2x_1 - 2x_2 \leq 2 \\ & 4x_1 - 3x_2 \leq 5 \\ & -x_1 + 2x_2 \leq 3 \\ & x_1 \geq 0, x_2 \geq 0 \end{aligned}$$

Turn the page for exercises 5 and 6!

8
8
8
b is a scalar, not a vector

5. Consider the following linear program in standard form (Phase 2):

$$\begin{aligned} \min \quad & \mathbf{c}'\mathbf{x} \\ \text{s.t.} \quad & \mathbf{Ax} = \mathbf{b} \\ & \mathbf{x} \geq \mathbf{0} \end{aligned}$$

Here matrix $\mathbf{A} \in \mathbb{R}^{m \times n}$ has m linearly independent rows, and vector $\mathbf{b} \in \mathbb{R}^m$, $\mathbf{b} \geq \mathbf{0}$. Recall that in the first phase of the two-phase simplex method, we introduce a vector $\mathbf{y} = (y_1, \dots, y_m)$, and solve the following linear program (Phase 1):

$$\begin{aligned} \min \quad & \sum_{i=1}^m y_i \\ \text{s.t.} \quad & \mathbf{Ax} + \mathbf{y} = \mathbf{b} \\ & \mathbf{x} \geq \mathbf{0} \\ & \mathbf{y} \geq \mathbf{0} \end{aligned}$$

- §
- §
- §
- §
- Explain how we can find a basic feasible solution to the Phase 1 problem.
 - Explain for which results of the Phase 1 problem, we can conclude that the Phase 2 problem is infeasible.
 - Explain for which results of the Phase 1 problem, we can conclude that the Phase 2 problem has a basic feasible solution, and explain how to find this basic feasible solution.
 - Instead of $\min \sum_{i=1}^m y_i$, we can also use a different objective in Phase 1. For both of the following two objectives, explain intuitively whether you can conclude from an optimal solution of Phase 1, whether the Phase 2 problem has a feasible solution or not.

Objective 1: $\min \sum_{i=1}^m i \cdot y_i$. Objective 2: $\min \sum_{i=1}^m -y_i$

6. Consider the following linear program with optimal solution

$$x_1 = 30/17, x_2 = -11/17, x_3 = 0.$$

$$\begin{aligned} \min \quad & 7x_1 - x_2 + 2x_3 \\ \text{s.t.} \quad & x_1 - 5x_2 + x_3 \geq 5 \\ & 3x_1 + 2x_2 + x_3 \geq 4 \\ & 2x_1 + 2x_2 - x_3 \geq 1 \\ & x_1 + 3x_2 \geq -1 \\ & x_1 \geq 0, x_2 \leq 0, x_3 \leq 0 \end{aligned}$$

- §
- §
- Construct the dual program.
 - Use complementary slackness to compute an optimal solution of the dual program.

exercise	1	2	3	4	5	6
points	8	8	8	8	8	8