

**201700080 Information Theory and Statistics**  
**8 April 2021, 9:00 – 12:00**

This test consists of 6 problems for a total of 26 points. All answers need to be justified. The use of a non-programmable calculator (not a “GR”) is allowed. A handwritten single side A4 cheat sheet is allowed. No additional books or notes may be used.

1. You receive a sample of  $n$  observations  $x_1, x_2, \dots, x_n$ , that are independent and identically distributed according to a Bernoulli distribution for which the success probability  $p$  is unknown. It is known that either  $p = p_1$  or  $p = p_2$ , i.e. the distribution for a single observation is either  $P_1(X = x) = (1 - p_1)^{1-x} p_1^x$  or  $P_2(X = x) = (1 - p_2)^{1-x} p_2^x$ ,  $x \in \{0, 1\}$ .
  - a. [3 pt] Specify a binary hypothesis testing problem for choosing between  $p_1$  and  $p_2$ . Derive an optimal decision rule.
  - b. [1 pt] Compute  $D(P_1 \parallel P_2)$ . (*Hint*:  $\mathbb{E}[X] = 1/p$ )
  - c. [2 pt] State the Chernoff-Stein result and discuss the implications for this problem.
2. Let  $X_1$  and  $X_2$  be identically distributed discrete random variables. They are not necessarily independent. Let

$$\rho = 1 - \frac{H(X_2|X_1)}{H(X_1)}.$$

- a. [2 pt] Show that  $\rho = \frac{I(X_1; X_2)}{H(X_1)}$ .
  - b. [1 pt] Show that  $0 \leq \rho \leq 1$ .
  - c. [1 pt] When is  $\rho = 0$ ?
  - d. [2 pt] When is  $\rho = 1$ ?
3. [3 pt] An engineer constructs a Huffman code that achieves expected codeword length  $\mathbb{E}[L] = 2.3$ . Give an upper and a lower bound on the Shannon entropy of the underlying source. Justify both bounds.
  4. [3 pt] A dictionary-based Lempel-Ziv code uses  $\lambda$  to denote the empty string and  $\sigma$  to denote the end-of-message symbol. It gives the following compressed message:

$(\lambda, a), (1, a), (\lambda, b), (3, c), (3, b), (2, a), (\lambda, \sigma)$ .

Decode and give the original message.

**P.T.O.** (Please turn over)

5. [3 pt] Consider the binary symmetric channel, i.e. the channel with  $\mathcal{X} = \mathcal{Y} = \{0, 1\}$  and  $P(Y = 0|X = 0) = P(X = 1|Y = 1) = 1 - p$  and  $P(Y = 1|X = 0) = P(X = 0|Y = 1) = p$ . Derive an expression for the capacity of this channel as a function of  $p$ .
6. Let  $X$  be discrete random variable for which  $p_\theta(k) = P(X = k) = \frac{\theta^k e^{-k}}{k!}$ , for  $k = 0, 1, 2, \dots$ . Here,  $k!$  denotes the factorial, i.e.  $0! = 1$  and  $k! = \prod_{i=1}^k i$  for  $k \geq 1$ . Random variables of this form satisfy  $\mathbb{E}[X] = \text{var}(X) = \theta$ .
- a. [3 pt] Assume that  $\theta$  is unknown and compute the Fisher information  $J(\theta)$ . (*Hint: For discrete probability distributions with parameter  $\theta$ , the Fisher information is given by  $J(\theta) = \mathbb{E}_\theta \left[ \left( \frac{\partial}{\partial \theta} \ln p_\theta(X) \right)^2 \right]$  .)*
- b. [2 pt] State the Cramer-Rao bound and discuss in your own words the implications of the result from a) on unbiased estimators of  $\theta$  based on  $n$  observations.